

# Symmetric Lenses

Martin Hofmann   Benjamin Pierce   Daniel Wagner

January 27, 2011  
POPL Austin

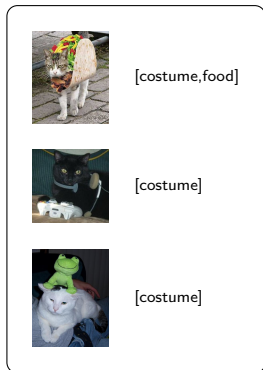
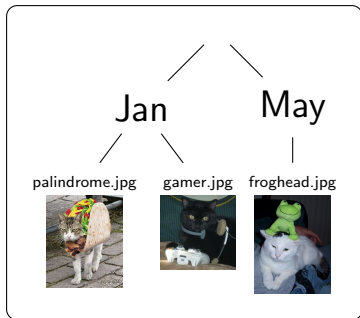
---

# Setup

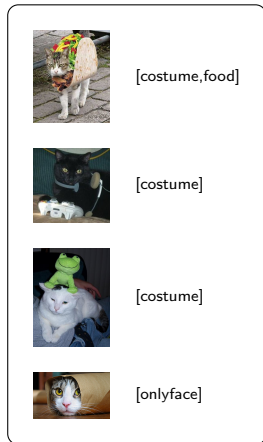
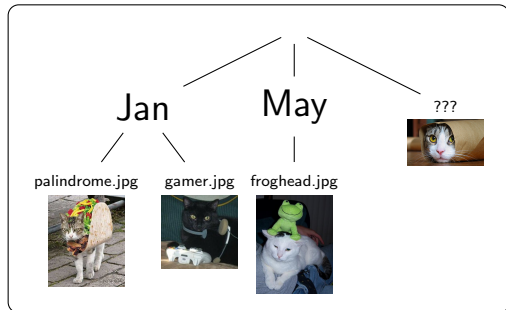
Daniel shares cat pictures with his coworkers, but prefers a different organization scheme than they do.

- ▶ At home: tree-structured file-system
  - ▶ On the web: flat-list picture gallery with tags
-

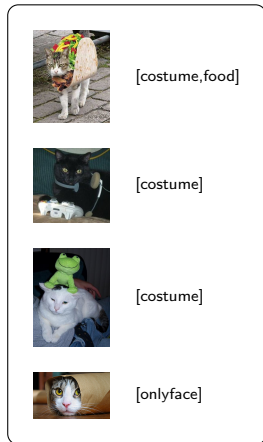
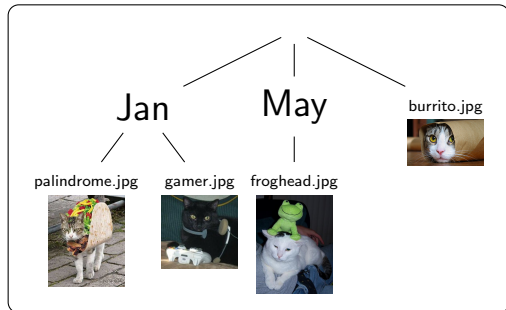
# Goal, Part 1: Two Structures



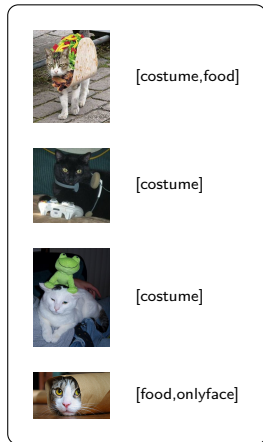
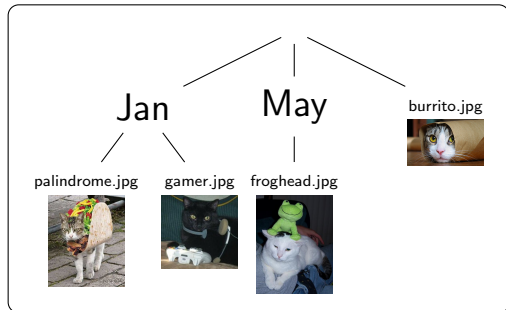
# Goal, Part 2: Adding to the Web Gallery



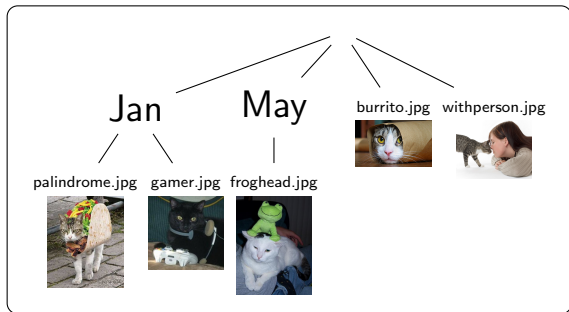
# Goal, Part 3: Fixing the Filename



# Goal, Part 4: Changing Tags



# Goal, Part 5: Adding to the File System



[costume,food]

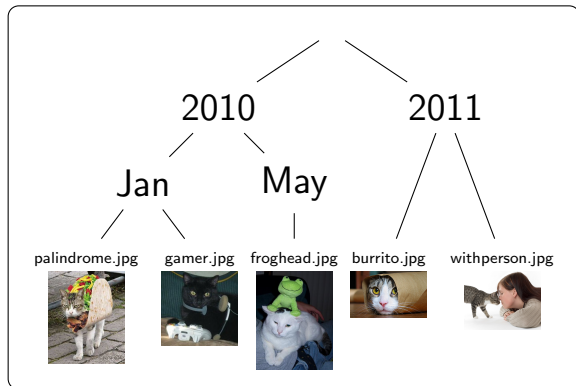
[costume]

[costume]

[food,onlyface]

[ ]

# Goal, Part 6: Restructuring



[costume,food]

[costume]

[costume]

[food,onlyface]

[ ]



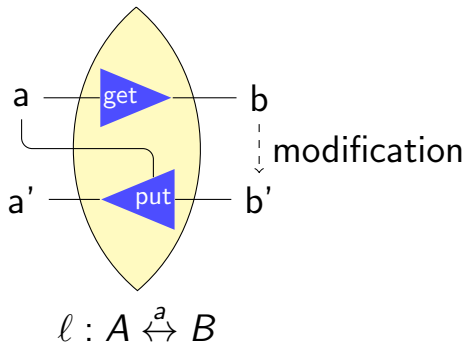
We Get It... Use Lenses

---

# Use Lenses, Dummy!

- ▶ Combinators for Bidirectional Tree Transformations  
(Foster, Greenwald, Moore, Pierce, Schmitt; POPL 2005)
  - ▶ Relational Lenses: A Language For Updateable Views  
(Bohannon, Vaughn, and Pierce; PODS 2006)
  - ▶ Boomerang: Resourceful Lenses for String Data  
(Bohannon, Foster, Pierce, Pilkiewicz, and Schmitt; POPL 2008)
  - ▶ Bidirectional Programming Languages  
(Foster; thesis 2009)
  - ▶ Bidirectionalizing Graph Transformations  
(Hidaka, Hu, Inaba, and Kato; ICFP 2010)
  - ▶ Update Semantics of Relational Views  
(Bancilhon and Spyratos; 1981)
-

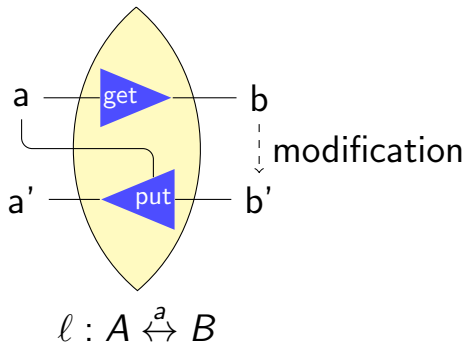
# Lenses 101



$\text{get} : A \rightarrow B$

$\text{put} : B \times A \rightarrow A$

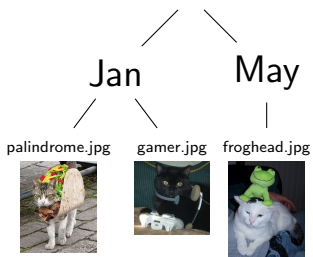
# Lenses 101



$get : A \rightarrow B$

$put : B \times A \rightarrow A$

# Typing “get”



[costume,food]



[costume]



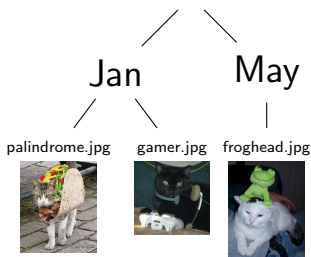
[costume]

data FS = Directory Name [FS]  
| File Name Picture

type Web =  
[(Picture, [Tag])]

$\ell : FS \overset{?}{\leftrightarrow} \text{Web}$

# Typing “get”



[costume,food]



[costume]



[costume]

data FS = Directory Name [FS]  
| File Name Picture

type Web =  
[(Picture, [Tag])]

$\ell : \text{Web} \stackrel{?}{\leftrightarrow} \text{FS}$

# Formalizing the Oddity: Roundtrip Laws

$$\textit{put}(\textit{get}(a), a) = a$$

$$\textit{get}(\textit{put}(b, a)) = b$$

Either possibility forbidden!

---

# Well, What About...

- ▶ Symmetric Constraint Maintainers  
(Meertens; 1998)
- ▶ Towards an Algebraic Theory of Bidirectional Transformations  
(Stevens; ICGT 2008)
- ▶ Bidirectional Model Transformations in QVT: Semantic Issues and Open Questions  
(Stevens; MoDELS 2007)
- ▶ Algebraic Models for Bidirectional Model Synchronization  
(Diskin; MoDELS 2008)
- ▶ Supporting Parallel Updates with Bidirectional Model Transformations  
(Xiong, Song, Hu, and Takeichi; ICMT 2009)

No composition!

---



# Our Contribution

A lens framework with

1. symmetry
  2. composition
-

# Our Contribution

A lens framework with

1. symmetry
  2. composition
  3. ... and other nice combinators
-

# Symmetrizing Lenses

---

# Starting Point: Asymmetric Lenses

$$l : A \overset{a}{\leftrightarrow} B$$

$$\text{get} : A \rightarrow B$$

$$\text{put} : B \times A \rightarrow A$$

$$\text{get}(\text{put}(b, a)) = b$$

$$\text{put}(\text{get}(a), a) = a$$

---

# L/R Symmetry

$$l : A \overset{a}{\leftrightarrow} B$$

$$\text{putr} : A \times B \rightarrow B$$

$$\text{putl} : B \times A \rightarrow A$$

$$\text{get}(\text{put}(b, a)) = b$$

$$\text{put}(\text{get}(a), a) = a$$

---

# Complements

$$l : A \overset{a}{\leftrightarrow} B$$

$$\text{putr} : A \times S_B \rightarrow B$$

$$\text{putl} : B \times S_A \rightarrow A$$

$$\text{get}(\text{put}(b, a)) = b$$

$$\text{put}(\text{get}(a), a) = a$$

---

# I/O Symmetry

$$l : A \overset{a}{\leftrightarrow} B$$

$$\text{putr} : A \times S_B \rightarrow B \times S_A$$

$$\text{putl} : B \times S_A \rightarrow A \times S_B$$

$$\text{get}(\text{put}(b, a)) = b$$

$$\text{put}(\text{get}(a), a) = a$$

---

# Unifying Complements

$$l : A \leftrightarrow B$$

$$\text{putr} : A \times S \rightarrow B \times S$$

$$\text{putl} : B \times S \rightarrow A \times S$$

$$\text{get}(\text{put}(b, a)) = b$$

$$\text{put}(\text{get}(a), a) = a$$

---



# Updated Lens Laws

$$l : A \leftrightarrow B$$

$$\text{putr} : A \times S \rightarrow B \times S$$

$$\text{putl} : B \times S \rightarrow A \times S$$

$$\frac{\text{putr}(a, s) = (b, s')}{\text{putl}(b, s') = (a, s)}$$

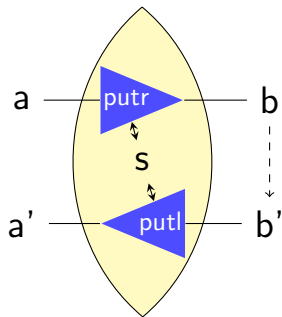
$$\frac{\text{putl}(b, s) = (a, s')}{\text{putr}(a, s') = (b, s)}$$

---

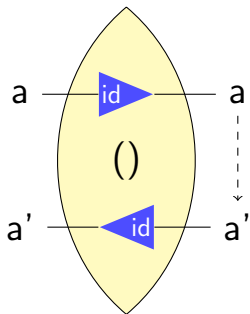
# Composition

---

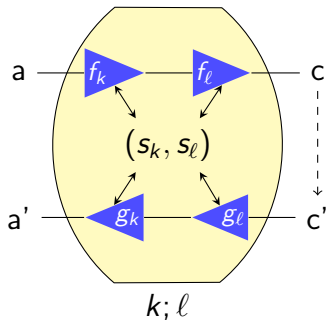
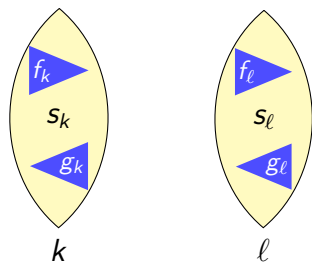
# Updated Wiring Diagram



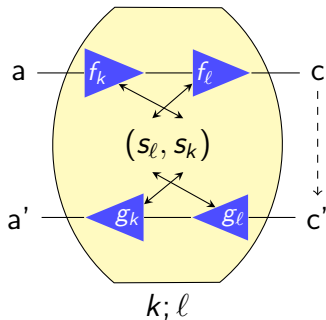
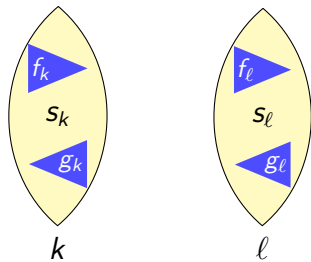
# Warm-up: Identity Lens



# Composition



# Another Composition



# Behavioral Equivalence

$k \equiv \ell$  when there's a relation  $R \subset k.S \times \ell.S$  and:

$$\frac{\begin{array}{l} s_k R s_\ell \\ k.putr(a, s_k) = (b_k, s'_k) \\ \ell.putr(a, s_\ell) = (b_\ell, s'_\ell) \end{array}}{b_k = b_\ell \wedge s'_k R s'_\ell}$$

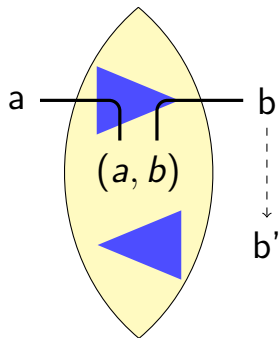
$$\frac{\begin{array}{l} s_k R s_\ell \\ k.putl(b, s_k) = (a_k, s'_k) \\ \ell.putl(b, s_\ell) = (a_\ell, s'_\ell) \end{array}}{a_k = a_\ell \wedge s'_k R s'_\ell}$$

# Handy Lenses

---

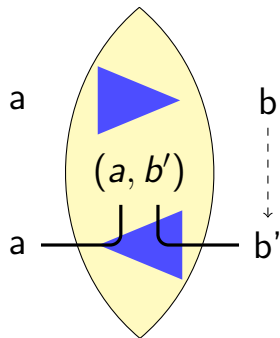


# Disconnect (putr)



disconnect :  $A \leftrightarrow B$

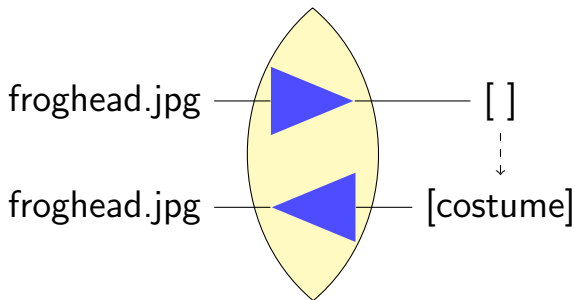
# Disconnect (putl)



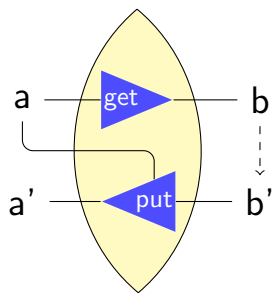
disconnect :  $A \leftrightarrow B$

---

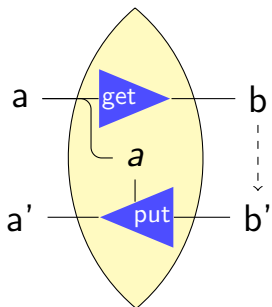
# Why disconnect?



# Lifting Asymmetric Lenses

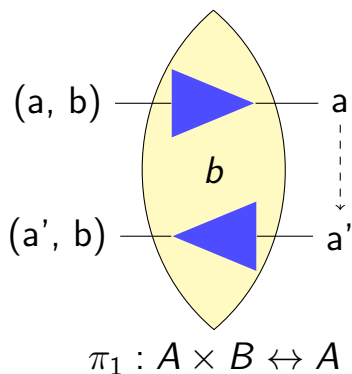


$$l : A \overset{a}{\leftrightarrow} B$$



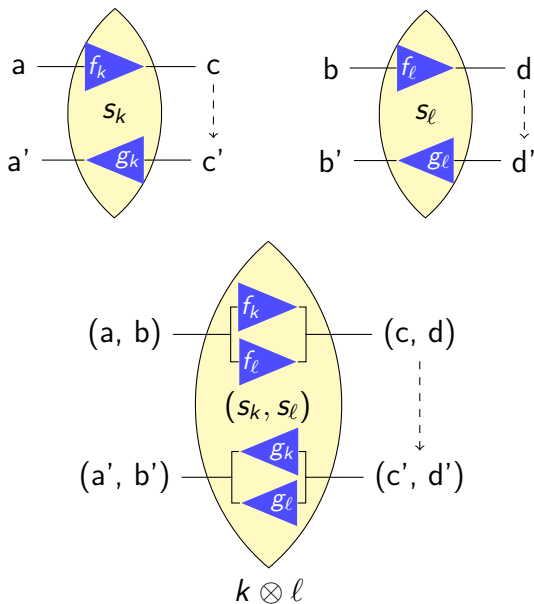
$$l^{sym} : A \leftrightarrow B$$

# Projection



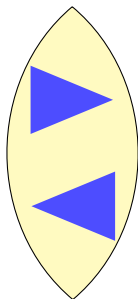
$(\pi_2 \text{ is similar})$

# Tensor Product



# Synchronizing Tree Leaves/List Elements

froghead.jpg

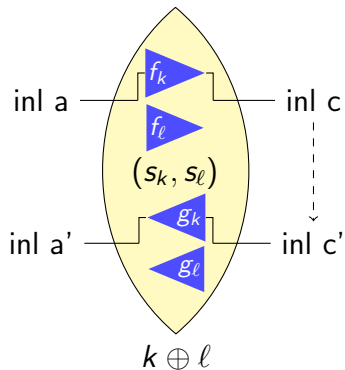
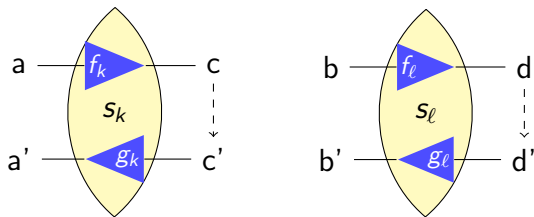


[costume]



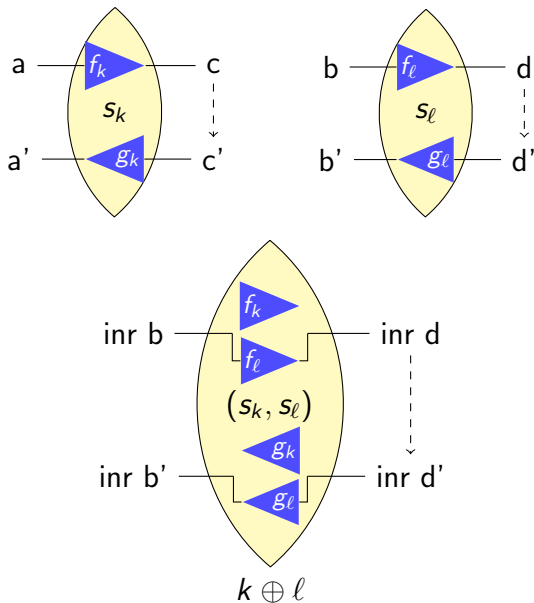
disconnect  $\otimes$  id

# Tensor Sum





# Tensor Sum



# Lenses for Recursive Data

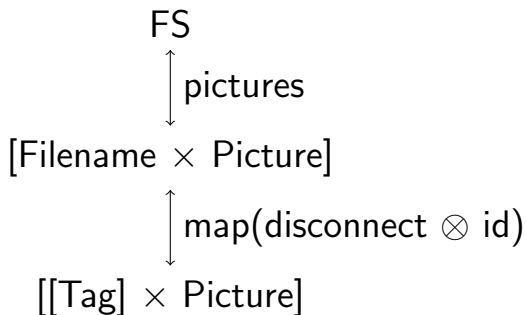
$$\triangleright \frac{\ell : F(A) \leftrightarrow A}{\text{fold}(\ell) : \mu F \leftrightarrow A}$$

- ▶ Initial algebra construction for data
  - ▶ Some technical requirements
  - ▶ Package up the fold and unfold operations
-

# Useful Folds

- ▶ leaves :  $\text{Tree } A \leftrightarrow [A]$
  - ▶ concat :  $[[A]] \leftrightarrow [A]$
  - ▶ partition :  $[A \uplus B] \leftrightarrow [A] \times [B]$
  - ▶ map :  $(A \leftrightarrow B) \rightarrow ([A] \leftrightarrow [B])$
  
  - ▶ pictures :  $\text{FS} \leftrightarrow [\text{Name} \times \text{Picture}]$
-

# Final Lens



# Conclusion

- ▶ Theoretical framework
    - ▶ Symmetric and compositional
    - ▶ Behavioral equivalence
  - ▶ Lens language
    - ▶ Miscellaneous useful basic lenses
    - ▶ Tensor sums and products, projections, injections
    - ▶ ADTs via folds and unfolds
    - ▶ Mapping for (non-algebraic) containers
  - ▶ Relationship to asymmetric lenses
    - ▶ Embedding of asymmetric lenses
    - ▶ Decomposition into asymmetric lens spans
-